Storage

*Database Systems: The Complete Book*
UBDB Seminars

Mondays @ 10:30 AM in Davis 113A

Feb 15: Rethinking the Database for the Data Science Era
        Zack Ives (UPenn)

Feb 22: Large-Scale Machine Learning With The SimSQL System
        Chris Jermaine (Rice)

Mar 21: Approximate lifted inference with probabilistic databases
        Wolfgang Gatterbauer (CMU)

April 18: Title TBD
         Ihab Ilyas (Waterloo)

http://odin.cse.buffalo.edu/seminar/index.html
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# Recap

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## Recap

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   & \quad <1,5> \quad <1,8> \} \\
2 & \quad \{ <2,3> \quad <2,7> \} \\
3 & \quad \{ <3,4> \quad <3,6> \}
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\end{align*}
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The Memory Hierarchy

Fast (but small)

Big (but slow)
Storage

• How do we...
  • ...optimize across the memory hierarchy?
  • ...use the right data access pattern for the storage medium we’re using?
  • ...organize data to minimize access costs?
  • ...organize data to minimize storage costs?
The IO Problem

Computations

READ

WRITE

Expensive
Why not use just RAM?

• RAM is more expensive than HD
  • 200 MB/$ vs 10 GB/$

• RAM is smaller
  • 128 GB vs 10 TB

• RAM is volatile

Are in-memory databases still viable? (Hint: Yes)
In-Memory DBs

- Why use In-Memory DBs?
  - Faster processing (especially for random access)
- How can we provide persistence?
  - … with respect to local failures (crashes)
  - … with respect to global failures (hurricanes)
- How do we provide scale?
  - Some DBs need TB/PB/EB of space.
def Select(predicate, source):
    while(source.hasMoreTuples):
        tuple = source.readTuple()
        if(predicate(tuple)):
            output(tuple)

High Latency if source is disk!

Where is output stored?
def Select(predicate, source)
    while(source.hasMoreTuples)
        in_buffer = source.fetch()
        while(in_buffer.hasMoreTuples)
            tuple = in_buffer.readTuple()
            if(predicate(tuple))
                out_buffer.output(tuple)
                if(out_buffer.isFull)
                    out_buffer.flush()
Data Organization

• How do we store data?
  • How are records represented on-disk? (Serialization)
  • How are records stored within a page?
  • How are pages organized in a file?
  • What other metadata do we need?
• Our solutions must also be persisted to disk.
Files and Data

- A **File** is a collection of pages
  - A **Page** is a collection of records
    - A **Record** is a data value (e.g., a tuple)
  - We need an infrastructure to ensure that records we need are in memory.
  - We need some way to organize and store files, pages, and records.
How is data laid out in a record?
Record (Tuple) Formats

What are some advantages/disadvantages of storing records this way?
Record (Tuple) Formats

Record information stored in a System Catalog

Base Address (B) Address $B + |L1| + |L2|$

What are some advantages/disadvantages of storing records this way?
Record (Tuple) Formats

Array of Field Offsets

What are some advantages/disadvantages of storing records this way?
How are records laid out in a page?
Page Formats

What are advantages/disadvantages of these formats?
Page Formats

Variable Size Records

What are advantages/disadvantages of this format?
Files of Records

File: A collection of pages of records that must support:

- Read a record
- Insert/Delete/Update a record
- Scan all records
Unordered (Heap) Files

Store records in no particular order

Disk pages are allocated/freed as file grows and shrinks

Support for record level operations by:
- Keeping track of pages in the file
- Keeping track of free space in each page
- Keeping track of records on each page

This data must be stored somewhere!
Unordered (Heap) Files

Each page contains 2 pointers plus data
Directories are a collection of pages (e.g., a linked list)
Directories point to all data pages
(entries can include # of free pages)
What are the advantages and disadvantages of each?
def Select(predicate, source):
    while(source.hasMoreTuples):
        in_buffer = source.fetch()
        while(in_buffer.hasMoreTuples):
            tuple = in_buffer.readTuple()
            if(predicate(tuple)):
                out_buffer.output(tuple)
                if(out_buffer.isFull):
                    out_buffer.flush()
IO + Buffering

Generalize & Standardize!

Have a component that handles buffering!
The Buffer Manager

API

Allocate a page
Deallocate a page
Read from a page
Write to a page

Database

Query Interpretation and Optimization
Relational Operators
Files/Data Accessors
Buffer Management
Disk Management
The Buffer Manager

Higher levels of the DB

Disk Page

Free Frame

Pages allocated to frames as per page replacement policy
Pinned Pages

• Pinning a page indicates that it is being used.
• The requestor must unpin the page when done.
  • The requestor must also indicate whether the page has been modified (with a ‘dirty’ bit)
• Dirty pages must be written to disk
• Pages may be requested multiple times
  • Use a pin count (reference count) to keep track.
• Concurrency Control/Recovery may require other operations when replacing a frame.
Buffer Replacement

- Frames are chosen for replacement by a buffer replacement policy.
- (e.g., LRU, MRU, Clock)
- Policy can have a big impact!
- Depends on the access pattern.
- What is a worst-case scenario for LRU?
  Hmmm… this sounds awfully familiar…
Hey… Oliver!

This sounds a lot like virtual memory!
Buffer Managers vs Virtual Memory

• Not a huge difference
  • Many lightweight DBs use VMem as a buffer manager!

• Reasons to implement an explicit buffer manager:
  • Control when and how paging happens.
    • e.g., better/more efficient prefetching.
  • Control what gets paged in/out.
    • e.g., better knowledge of data access patterns.
  • Integrate additional memory layers (e.g., Network)
Example-OS Paging

Aha, a Sequential Read!

Index: ‘Enterprise officers in this range’
Example-DB Paging

Read

Read in precisely what you need.

Index: ‘Enterprise officers in this range’
Time permitting…
The record is the main unit of computation...

... but what if the records are really really really really big
CREATE TABLE visitor(
    id big_int,
    ip int,
    age int,
    gender enum,
    ...
    region string,
    country string,
    city string,
    ...
    likes_cats bool,
    likes_spring_break bool,
    likes_cookies bool,
    ...
);

Google, Facebook, Amazon, etc... have log files and customer information tables with 100s or 1,000s of columns.
Visitors

Tid

σ\text{likes} \cdot \text{season} = \text{spring-break}